



**CS606- compiler instruction**  
**Solved MCQS**  
**From Midterm Papers**

**Feb 25,2013**

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**PSMD01**

**Final Term MCQ's and Quizzes**  
**CS606- compiler instruction**

**Question No: 1 ( Marks: 1 ) - Please choose one**

If X is a terminal in  $A \rightarrow aX\cdot$ , then this transition corresponds to a shift of \_\_\_\_ from input to top of parse stack.

- X
- A
- a
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

A canonical collection of sets of items for an augmented grammar, C is constructed as -----

The first set in C is the closure of  $\{[S' \rightarrow \cdot S]\}$ , where S is starting symbol of original grammar and S' is the starting non-terminal of augmented grammar.

The first set in C is the closure of  $\{[S' \rightarrow \cdot S]\}$ , where S is starting symbol of original grammar and S' is the starting non-terminal of original grammar.

The first set in C is the closure of  $\{[S' \rightarrow \cdot S]\}$ , where S is starting symbol of original grammar and S is the starting non-terminal of augmented grammar.

None of these

**Question No: 1 ( Marks: 1 ) - Please choose one**

An ----- does not need to examine the entire stack for a handle, the state symbol on the top of the stack contains all the information it needs.

- LR parser
- RL parser
- BU parser
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Suppose ? begins with symbol X which may be a terminal (token) or non-terminal. The item can be written as A? Xa•?.

- True
- False

**Question No: 1 ( Marks: 1 ) - Please choose one**

YACC parser generator builds up

- SLR parsing table
- Canonical LR parsing table
- LALR parsing table
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

LR(1) parsing is --- base parsing.

- DFA
- CFG
- PDA
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

The LR(1) parsers can not recognize precisely those languages in which one-symbol lookahead suffices to determine whether to shift or reduce.

- True
- False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Yacc contains built-in support for handling ambiguous grammars resulting in shift-reduce conflicts. By default these conflicts are solved by performing the \_\_\_\_\_.

- Shift action
- Reduce action
- Shift and reduce actions
- De-allocation of memory

**Question No: 1 ( Marks: 1 ) - Please choose one**

$S \rightarrow A \mid xB$   $A \rightarrow aAb \mid x$  This grammar contains a reduce-reduce conflict.

True

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

$S \rightarrow a \mid B$

$B \rightarrow Bb \mid E$  The non-terminal \_\_\_\_\_ is left recursive.

B

a

E

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Following statement represents: if x relop y goto L

abstract jump

Conditional jump

While loop

None of the Given

**Question No: 1 ( Marks: 1 ) - Please choose one**

When generating a lexical analyzer from a \_\_\_\_\_ description, the item sets (states) are constructed by two types of “moves”: character moves and e moves.

Character

Grammar

Token

Sentence

**Question No: 1 ( Marks: 1 ) - Please choose one**

Left factoring is enough to make a grammar LL(1).

True

False



**Question No: 1 ( Marks: 1 ) - Please choose one**

Register allocation by graph coloring uses a register interference graph. \_\_\_\_\_ nodes in the graph are joined by an edge when the live ranges of the values they represent overlap.

**Two p116**

Three

Four

Five

**Question No: 1 ( Marks: 1 ) - Please choose one**

$S \rightarrow ABA \rightarrow e \mid aAB \rightarrow e \mid bB$  - FIRST(S) contains \_\_\_\_ elements.

**3**

4

5

6

**Question No: 1 ( Marks: 1 ) - Please choose one**

The notation \_\_\_\_\_ instructs YACC to push a computed attribute value on the stack.

**\$\$ Page no : 98**

&&

##

--

**Question No: 1 ( Marks: 1 ) - Please choose one**

Simple code generation considers one AST node at a time. If the target is a register machine, the code can be generated in one \_\_\_\_\_ traversal of the AST, possibly introducing temporaries when running out of registers.

Depth-first

Breadth-first

Top-Down

Bottom-Up

**Question No: 1 ( Marks: 1 ) - Please choose one**

Grammars with LL(1) conflicts can be made LL(1) by applying left-factoring, substitution, and left-recursion removal. Left-factoring takes care of \_\_\_\_\_conflicts.

FIRST/FIRST

FIRST/SECOND

SECOND/FIRST

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In an attribute grammar each production rule( $N \rightarrow a$ ) has a corresponding attribute evaluation rule that describes how to compute the values of the \_\_\_\_\_attributes of each particular node N in the AST.

Synthesized page no : 92

Complete

Free

Bounded

**Question No: 1 ( Marks: 1 ) - Please choose one**

When constructing an LR(1) parser we record for each item exactly in which context it appears, which resolves many conflicts present in \_\_\_\_\_parsers based on FOLLOW sets.

SLR(1)

LRS(1)

RLS(1)

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

The \_\_\_\_\_translation statements can be conveniently specified in YACC

Syntax-directed Page no : 120

Image-directed

Sign-directed

None of the given.

**Question No: 1 ( Marks: 1 ) - Please choose one**

Backpatching to translate flow-of-control statements in \_\_\_\_ pass.

- one
- two
- three
- all of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Alternative of the backtrack in parser is Look ahead symbol in \_\_\_\_\_ .

Input

- Output
- Input and Output
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Typical compilation means programs written in high-level languages to low-level \_\_\_\_\_.

Object code

- Byted code
- Unicode
- Both Object Code and byte code

**Question No: 1 ( Marks: 1 ) - Please choose one**

In PASCAL \_\_\_\_\_ represent the inequality test.

- :
- :=
- =



None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

LR parsing \_\_\_\_\_ a string to the start symbol by inverting productions.

Reduces

- Shifts
- Adds
- None of the given



**Question No: 1 ( Marks: 1 ) - Please choose one**

In multi pass compiler during the first pass it gathers information about \_\_\_\_\_ .

Declaration

Bindings

Static information

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ phase which supports macro substitution and conditional compilation.

**Semantic**

Syntax

Preprocessing

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In parser the two LL stand(s) for \_\_\_\_\_ .

Left-to-right scan of input

left-most derivation

**All of the given**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Parser always gives a tree like structure as output

**True**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Lexer and scanner are two different phases of compiler

True

**False**

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ tree in which each node represents an operator and children of the node represent the operands.

**Abstract syntax**    **Page no : 100**

Concrete syntax

Parse

None of the given

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**Question No: 1 ( Marks: 1 ) - Please choose one**

In compilation process Hierarchical analysis is also called

Parsing

**Syntax analysis**

Both Parsing and Syntax analysis

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Ambiguity can easily be handled by Top-down Parser

Select correct option:

**True**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Front-end of a two pass compiler is consists of Scanner.

**True**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

LL(1) parsing is called non-predictive parsing.

**True**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

In predictive parsing table the rows are \_\_\_\_\_ .

**Non-terminals**

Terminals

Both non-terminal and terminals

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In LL1() parsing algorithm \_\_\_\_\_ contains a sequence of grammar symbols.

**Stack**

Link list

Array

None



**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the grammar

$A \rightarrow B C D$

$B \rightarrow h B \mid \epsilon$

$C \rightarrow C g \mid g \mid C h \mid i$

$D \rightarrow A B \mid \epsilon$

First of C is \_\_\_\_\_ .

Select correct option:

g, I look down for reference

g

h i

i

**Question No: 1 ( Marks: 1 ) - Please choose one**

AST summarizes the grammatical structure with the details of derivations.

True

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Left factoring is enough to make LL1 grammar

True

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

If X is a non-terminal in A?  $aX\bullet$ ?, then the interpretation of this transition is more complex because non-terminals do not appear in input

Yes

No

**Question No: 1 ( Marks: 1 ) - Please choose one**

If / is a set of items for a grammar then closure (/) is a set of items constructed from / by the following rule.

If  $A \rightarrow aX.Y$  is in closure (/) and  $Y \rightarrow r$  is production, then add  $X \rightarrow .r$  to closure (/).

If  $A \rightarrow a.XY$  is in closure (/) and  $X \rightarrow r$  is production, then add  $X \rightarrow .r$  to closure (/).

If  $A \rightarrow aXY.$  is in closure (/) and  $A \rightarrow r$  is production, then add  $X \rightarrow .r$  to closure (/).

None of these

**Question No: 1 ( Marks: 1 ) - Please choose one**

NFA of LR(0) items means \_\_\_\_\_

- look-ahead one sybole
- no look-ahead
- look-ahead all sybols
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

A grammar is LR if a ----- shift reduce-reduce parser can recognize handles when they appear on the top of stack.

- left-to-reverse
- left-to-rise
- left-to-right
- None of the given.

**Question No: 1 ( Marks: 1 ) - Please choose one**

The output from the algorithm of constructing the collection of canonical sets of LR(1) items will be the \_\_\_\_\_

- Original Grammar G
- Augmented grammar G'
- Parsing table
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Reduction of a handle to the ----- on the left hand side of the grammar rule is a step along the reverse of a right most derivation.

- Terminal
- Non-terminal

**Question No: 1 ( Marks: 1 ) - Please choose one**

NFA of LR(1) items means \_\_\_\_\_

- no look-ahead
- look-ahead one sybole
- look-ahead all sybols
- None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In canonical collection procedure a DFA can not be constructed from NFA using the subset construction, similar to one we used for lexical analysis.

True  
False

**Question No: 1 ( Marks: 1 ) - Please choose one**

performing common subexpression elimination on aa dependency graph requires the identification of nodes with the same operator and operands. when using a hash table (with a hash function based on operator and operands) all \_\_\_\_\_ nodes can be identified in linear time.

common  
uncommon  
next  
previous

**Question No: 1 ( Marks: 1 ) - Please choose one**

Linear IRs resemble pseudo-code for same \_\_\_\_\_.

Automated Machine  
Mechanical machines  
Token machines  
Abstract machine



**Question No: 1 ( Marks: 1 ) - Please choose one**

The regular expressions  $a^*|b^*$  and  $(a|b)^*$  describe the \_\_\_\_\_ set of strings.

Same

**Different**

Onto

**Question No: 1 ( Marks: 1 ) - Please choose one**

Back patching to translate flow-of-control statements in \_\_\_\_\_ pass.

**one Page no : 111**

two

three

all of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the following grammar,  $S \rightarrow aT Ue$   $T \rightarrow Tbc/b$   $U \rightarrow d$  And suppose that string “abbcd e” can be parsed bottom-up by the following reduction steps: (i)  $aTbcde$  (ii)  $aTde$  (iii)  $aT Ue$  (iv)  $S$  So what can be a handle from the following?

The second (b) in (abbcd e)

The first (b) in (abbcd e)

The substring (cd) in (abbcd e)

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Yacc contains built-in support for handling ambiguous grammars resulting in \_\_\_\_\_ conflicts.

**Shift-reduce**

Shift-Shift

Shift-second

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

A lexical analyzer generator automatically constructs a \_\_\_\_\_ that recognizes tokens.

:

**FA**

PDA

DP

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Attributes whose values are defined in terms of a node's own attributes, node's siblings and node's parent are called \_\_\_\_\_ .

**Inherited attributes Page no : 92**

Physical attributes

Logical attributes

Un-synthesized attributes

**Question No: 1 ( Marks: 1 ) - Please choose one**

The following two items  $A \rightarrow P \bullet Q$   $B \rightarrow P \bullet Q$  can co-exist in an \_\_\_\_\_ item set.

LR

LS

LT

PR

**Question No: 1 ( Marks: 1 ) - Please choose one**

Three-address codes are often implemented as a \_\_\_\_\_.

**Set of quadruples Page no : 104**

Set of doubles

Set of Singles

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

The error handling mechanism of the yacc parser generator pushes the input stream back when inserting 'missing' tokens.

True

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Flow of values used to calculate synthesized attributes in the parse tree is:

**Bottom-up Page no: 92**

Right to left

Top-Down

Left to right

**Question No: 1 ( Marks: 1 ) - Please choose one**

What does following statement represent?  $x[i] = y$

Prefix assignment

Postfix assignment

**indexed assignment**      **Page no : 107**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

A lexical analyzer transforms a stream of tokens. The tokens are stored into symbol table for further processing by the parser.

**True Page no: 99**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

LR parsers can handle \_\_\_\_\_ grammars.

**Left-recursive Page no: 163**

file-recursive

End-recursive

Start-recursive

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ convert the reloadable machine code into absolute machine code by linking library and reloadable object files.

Assembler

**Loader/link-editor**

Compiler

Preprocessor

**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the following grammar,

$A \rightarrow B C D$

$B \rightarrow h B \mid \text{episilon}$

$C \rightarrow C g \mid g \mid C h \mid i$

$D \rightarrow A B \mid \text{episilon}$

First of A is \_\_\_\_\_ .

**h, g, i**

gh

None of the given

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**Question No: 1 ( Marks: 1 ) - Please choose one**

One of the core tasks of compiler is to generate fast and compact executable code.

**True**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Compilers are sometimes classified as.

Single pass

Multi pass

Load and go

**All of the given**

**Question No: 1 ( Marks: 1 ) - Please choose one**

In multi pass compiler during the first pass it gathers information about \_\_\_\_\_ .

Declaration

Bindings

Static information

**None of the given \*\***

**Question No: 1 ( Marks: 1 ) - Please choose one**

For each language to make LL(1) grammar, we take two steps, 1st is removing left recurrence and 2nd is applying fin sequence.

True

**False**

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ is evaluated to yield a value.

Command

**Expression**

Declaration

Declaration and Command

**Question No: 1 ( Marks: 1 ) - Please choose one**

We can get an LL(1) grammar by \_\_\_\_\_ .

Removing left recurrence

Applying left factoring

**Removing left recurrence and Applying left factoring**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Can a DFA simulate NFA?

Yes

No

Sometimes

Depend upon nfa

**Question No: 1 ( Marks: 1 ) - Please choose one**

Which of the statement is true about Regular Languages?

Regular Languages are the most popular for specifying tokens.

Regular Languages are based on simple and useful theory.

Regular Languages are easy to understand.

All of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

The transition graph for an NFA that recognizes the language  $(a \mid b)^*abb$  will have following set of states.

{0}

{0,1}

{0,1,2}

{0,1,2,3} not sure

**Question No: 1 ( Marks: 1 ) - Please choose one**

Functions of Lexical analyzer are?

Removing white space

Removing constants, identifiers and keywords

Removing comments

All of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the following grammar,  $S \rightarrow aT Ue$   $T \rightarrow Tbc/b$   $U \rightarrow d$  And suppose that string "abbcd" can be parsed bottom-up by the following reduction steps: (i)  $aTbcde$  (ii)  $aTde$  (iii)  $aT Ue$  (iv)  $S$  So, what can be a handle from the following?

The whole string, (aT Ue)      Page no : 68

The whole string, (aTbcde)

The whole string, (aTde)

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

The LR(1) items are used as the states of a finite automaton (FA) that maintains information about the parsing stack and progress of a shift-reduce parser.

**True Page no: 74**

**False**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Flex is an automated tool that is used to get the minimized DFA (scanner).

True

**False Page no: 26**

**Question No: 1 ( Marks: 1 ) - Please choose one**

We use ----- to mark the bottom of the stack and also the right end of the input when considering the Stack implementation of Shift-Reduce Parsing.

Epsilon

#

**\$ Page no : 65**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

When generating a lexical analyzer from a token description, the item sets (states) are constructed by two types of “moves”: character moves and \_\_\_\_ moves.

**E (empty string) Page no : 18**

#

@

none of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Bottom-up parsers handle a \_\_\_\_\_ class of grammars.

**large Page no : 63**

small

medium

none of the given



**Question No: 1 ( Marks: 1 ) - Please choose one**

Let a grammar  $G = (V_n, V_t, P, S)$  is modified by adding a unit production  $S' \rightarrow S$  to the grammar and now starting non-terminals becomes  $S'$  and grammar becomes  $G' = (V_n \cup \{S'\}, V_t, P \cup \{S' \rightarrow S\}, S')$ . The Grammar  $G'$  is called the -----

**Augmented Grammar Page no : 76**

Lesser Grammar

Anonymous Grammar

none of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Parser takes tokens from scanner and tries to generate \_\_\_\_\_ .

Binary Search tree

Parse tree

**Syntax trace Page no : 6**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In Flex specification file different sections are separated by \_\_\_\_\_ .

**%% Page no: 26**

&&

##

\\

**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the grammar  $A \rightarrow B C D$

$B \rightarrow h B \mid \epsilon$

$C \rightarrow C g \mid g \mid C h \mid i$

$D \rightarrow A B \mid \epsilon$

Follow of B is \_\_\_\_\_ .

h

g, h, i, \$

g, i

g

**Question No: 1 ( Marks: 1 ) - Please choose one**

Consider the grammar  $A \rightarrow B C D$

$B \rightarrow h B \mid \epsilon$

$C \rightarrow C g \mid g \mid C h \mid i$

$D \rightarrow A B \mid \epsilon$

Follow of C is \_\_\_\_\_ .

**g, h, i, \$ Page no : 47**

g, h, \$

h, i, \$

h, g, \$

**Question No: 1 ( Marks: 1 ) - Please choose one**

In DFA minimization we construct one \_\_\_\_\_ for each group of states from the initial DFA.

**State Page no : 25**

NFA

PDA

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

An important component of semantic analysis is \_\_\_\_\_ .

code checking

**type checking page no : 6**

flush checking

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Intermediate Representation (IR) stores the value of its operand in \_\_\_\_\_ .

**Registers Page no : 10**

Memory

Hard disk

Secondary storage

**Question No: 1 ( Marks: 1 ) - Please choose one**

In \_\_\_\_\_ certain checks are performed to ensure that components of a program fit together meaningfully.

Linear analysis

Hierarchical analysis

**Semantic analysis Page no : 33**

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Which of the following statement is true about Two pass compiler.

Front End depends upon Back End

**Back End depends upon Frond End**      **page no : 5**

Both are independent of each other

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ algorithm is used in DFA minimization.

James's

Robert's

**Hopcroft's**      **Page no : 19**

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

A \_\_\_\_\_ is a top down parser.

**Predictive Parsing** **Page no: 46**

Reactive parser

Proactive parser

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Lexical Analyzer generator \_\_\_\_\_ is written in Java.

Flex

**Jlex** **Page no : 26**

Complex

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ avoid hardware stalls and interlocks.

Register allocation

**Instruction scheduling** **Page no : 10**

Instruction selection

None of given



**Question No: 1 ( Marks: 1 ) - Please choose one**

Recursive \_\_\_\_\_ parsing is done for LL(1) grammar.

**Decent Page no : 47**

Ascent

Forward

Backward

**Question No: 1 ( Marks: 1 ) - Please choose one**

Left factoring of a grammar is done to save the parser from back tracking.

**True Page no:61**

False

**Question No: 1 ( Marks: 1 ) - Please choose one**

Responsibility of \_\_\_\_\_ is to produce fast and compact code.

**Instruction selection**

Register allocation

Instruction scheduling

**None of given Page no: 9**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Optimal registers allocation is an NP-hard problem.

True

**False Page no : 10**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Front end of two pass compiler takes \_\_\_\_\_ as input.

**Source code Page no: 5**

Intermediate Representation (IR)

Machine Code

None of the Given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In Three-pass compiler \_\_\_\_\_ is used for code improvement or optimization.

Front End

**Middle End Page no : 10**

Back End

Both Front end and Back end

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**Question No: 1 ( Marks: 1 ) - Please choose one**

\_\_\_\_\_ of a two-pass compiler is consists of Instruction selection, Register allocation and Instruction scheduling.

**Back end Page no : 9**

Front end

Start

None of given

**Question No: 1 ( Marks: 1 ) - Please choose one**

NFA is easy to implement as compared to DFA.

True

**False Page no : 19**

**Question No: 1 ( Marks: 1 ) - Please choose one**

In Back End module of compiler, optimal register allocation uses\_\_\_\_\_.

$O(\log n)$

$O(n \log n)$

**N P-Complete Page no : 10**

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

In a transition table cells of the table contain the \_\_\_\_\_ state.

Reject state

**Next state Page no 18**

Previous state

None of the given

**Question No: 1 ( Marks: 1 ) - Please choose one**

Parser generator for the grammar LALR (1) is:

**YACC, Bison, CUP Page no: 88**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Attributes of a node whose values are defined wholly in terms of attributes of node's children and from constants are called \_\_\_\_\_.

**Synthesized attributes Page no : 92**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Goto L statement represent

**Unconditional jump Page no : 107**

**Question No: 1 ( Marks: 1 ) - Please choose one**

Dotted items ( $T \square a \bullet b$ ) record which part of a token has already been matched. Integer?  $([0-9])^+$  • This is a \_\_\_\_\_ item.

**Extended Page no : 73**

**Question No: 1 ( Marks: 1 ) - Please choose one**

If  $T \rightarrow XYZ$  is a production of grammar  $G$  then which of the following item indicates that a string derivable from  $X$  has been seen so far on the input and we hope to see a string derivable from  $YZ$  next on the input.

**Question No: 1 ( Marks: 1 ) - Please choose one**

The most powerful parser is:

**Question No: 1 ( Marks: 1 ) - Please choose one**

In the Parsing Table the rows correspond to Parsing DFA states and columns correspond to ----.

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